## REPORT DOCUMENTATION PAGE Form Approved OMB NO. 0704-0188 The public reporting burden for this collection of information is estimated to average 1 hour per response, including the time for reviewing searching existing data sources, gathering and maintaining the data needed, and completing and reviewing the collection of information. regarding this burden estimate or any other aspect of this collection of information, including suggesstions for reducing this burden, to Washington Headquarters Services, Directorate for Information Operations and Reports, 1215 Jefferson Davis Highway, Suite 1204, Arlington VA, 22202-4302. Respondents should be aware that notwithstanding any other provision of law, no person shall be subject to any oenalty for failing to comply with a collection of information if it does not display a currently valid OMB control number. PLEASE DO NOT RETURN YOUR FORM TO THE ABOVE ADDRESS. 2. REPORT TYPE 1. REPORT DATE (DD-MM-YYYY) 3. DATES COVERED (From - To) 25-Jul-2009 - 31-Aug-2012 18-12-2012 Final Report 4. TITLE AND SUBTITLE 5a. CONTRACT NUMBER Final Report W911NF-09-1-0391 5b. GRANT NUMBER 5c. PROGRAM ELEMENT NUMBER 611102 6. AUTHORS 5d. PROJECT NUMBER A. P. Young 5e. TASK NUMBER 5f. WORK UNIT NUMBER 7. PERFORMING ORGANIZATION NAMES AND ADDRESSES 8. PERFORMING ORGANIZATION REPORT NUMBER University of California - Santa Cruz Office of Sponsored Projects The Regents of the University of California, Santa Cruz Santa Cruz, CA 95064 -1077 9. SPONSORING/MONITORING AGENCY NAME(S) AND 10. SPONSOR/MONITOR'S ACRONYM(S) ADDRESS(ES) ARO 11. SPONSOR/MONITOR'S REPORT U.S. Army Research Office NUMBER(S) P.O. Box 12211 Research Triangle Park, NC 27709-2211 56290-PH-OC.7 12. DISTRIBUTION AVAILIBILITY STATEMENT Approved for Public Release; Distribution Unlimited 13. SUPPLEMENTARY NOTES The views, opinions and/or findings contained in this report are those of the author(s) and should not contrued as an official Department of the Army position, policy or decision, unless so designated by other documentation.

### 14. ABSTRACT

Quantum Monte Carlo simulations, supplemented by diagonalization, were used to see how efficiently a quantum computer could solve optimization problems using the quantum adiabatic algorithm (QAA). Comparisons were made with a classical heuristic algorithm, WalkSAT. A preliminary study was also made to see if the QAA could solve the graph isomorphism problem. Several optimization problems were considered. Those of the K-SAT type, which, in mean field theory have a "random first order transition", were found to have an exponentially small

### 15. SUBJECT TERMS

Quantum Adiabatic Algorithm, Optimization, Monte Carlo, quantum computer, satisfiability problems, spin glass, quantum phase transition

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a. REPORT b. ABSTRACT c. THIS PAGE		ABSTRACT	OF PAGES	Peter Young	
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					831-459-4151

## Final Report

### **ABSTRACT**

Quantum Monte Carlo simulations, supplemented by diagonalization, were used to see how efficiently a quantum computer could solve optimization problems using the quantum adiabatic algorithm (QAA). Comparisons were made with a classical heuristic algorithm, WalkSAT. A preliminary study was also made to see if the QAA could solve the graph isomorphism problem. Several optimization problems were considered. Those of the K-SAT type, which, in mean field theory have a "random first order transition", were found to have an exponentially small energy gap at the quantum phase transition. We also studied a spin glass problem in which the gap was found to be exponentially small not at the phase transition but beyond it in the "quantum spin glass" phase. These results show that that the simplest implementation of the QAA would require an amount of computer time which increases exponentially with problem size. For the future, it would be useful to test whether better results could be obtained by considering cleverer paths in "Hamiltonian space", and by considering different types of problem where the absolute ground state is not needed, only a good approximation to it.

Enter List of papers submitted or published that acknowledge ARO support from the start of the project to the date of this printing. List the papers, including journal references, in the following categories:

(a) Papers published in peer-reviewed journals (N/A for none)

TOTAL:	4
12/06/2012	5.00 Edward Farhi, David Gosset, Itay Hen, A. Sandvik, Peter Shor, A. Young, Francesco Zamponi. Performance of the quantum adiabatic algorithm on random instances of two optimization problems on regular hypergraphs, Physical Review A, (11 2012): 0. doi: 10.1103/PhysRevA.86.052334
12/06/2012	3.00 A. P. Young, Itay Hen. Exponential complexity of the quantum adiabatic algorithm for certain satisfiability problems, Physical Review E, (12 2011): 0. doi: 10.1103/PhysRevE.84.061152
12/06/2012	2.00 A. P. Young, V. N. Smelyanskiy. First-Order Phase Transition in the Quantum Adiabatic Algorithm, Physical Review Letters, (01 2010): 0. doi: 10.1103/PhysRevLett.104.020502
12/06/2012	4.00 Itay Hen, A. Young. Solving the graph-isomorphism problem with a quantum annealer, Physical Review A, (10 2012): 0. doi: 10.1103/PhysRevA.86.042310
	<u>Paper</u>
Received	Paner

Number of Papers published in peer-reviewed journals:

(b) Papers published in non-peer-reviewed journals (N/A for none)

Received	<u>Paper</u>				
08/17/2011	classical algo	etti, A. Young. Complexity of sevorithm WalkSAT, view E, (7 2011): 0. doi: 10.110		on problems using the he	euristic
TOTAL:	1				
Number of Pap	ers published in nor	n peer-reviewed journals:			
		(c)	Presentations		
Colloquium on	``Complexity of the O	Quantum Adiabatic Algorithm", Pr	inceton, September 24, 200	9.	
	on ``Complexity of thinin, Taiwan, Decem	ne Quantum Adiabatic Algorithm" a ber 1922, 2009	at Winter School on Quantu	m Information Science at N	Natl.~Cheng
Talk on ``First o	order phase transition	s in the Quantum Adiabatic Algori	thm" at the March meeting	of the APS, March 1519,	2010.
	•	nsition in the Quantum Adiabatic App., Trieste, June 14-18, 2010	Algorithm", at the Conferen	ce on Quantum Statistical M	Mechanics,
Talk on ``First o	order phase transition	in the Quantum Adiabatic Algorit	hm", at Statphys24, Cairns,	Australia, July 19-23.	
	•	antum Adiabatic Algorithm" at the Lyo, February 1416, 2011.	third international worksho	op on dynamics and manipu	llation of
Invited talk on " Spain, July 11		Algorithms" at Summer School on	"Quantum Information med	ets Statistical Mechanics" at	t El Escorial,
		lgorithms" at the International World (1), J\"ulich Supercomputing Cen	-	-	ystems for
	`Mind the gap, solvin	ng optimization problems with a qu 8, 2012.	nantum computer", NASA (	Quantum Future Technologi	ies
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	Non Po	eer-Reviewed Conference P	roceeding publication	s (other than abstract	ts):

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TOTAL:

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Number of No	on Peer-Reviewed Conference Proceeding publications (other than abstracts):			
	Peer-Reviewed Conference Proceeding publications (other than abstracts):			
Received	<u>Paper</u>			
12/06/2012	6.00 A.P. Young, S. Knysh, V.N. Smelyanskiy. The complexity of the Quantum Adiabatic Algorithm, Conference on Computational Physics Kaohsiung, Taiwan, . , . : ,			
TOTAL:	1			
Number of Pe	eer-Reviewed Conference Proceeding publications (other than abstracts):			
	(d) Manuscripts			
Received	<u>Paper</u>			
TOTAL:				
Number of Manuscripts:				
	Books			
Received	<u>Paper</u>			

TOTAL:

**Patents Submitted** 

**Patents Awarded** 

Awards

Elected member of the American Academy of Arts and Sciences in 2012

Awarded the 2009 Aneesur Rahman Prize for Computational Physics from the American Physical Society.

### **Graduate Students**

NAME PERCENT\_SUPPORTED

FTE Equivalent: Total Number:

### **Names of Post Doctorates**

<u>NAME</u>	PERCENT SUPPORTED
Itay Hen	1.00
FTE Equivalent:	1.00
Total Number:	1

## Names of Faculty Supported

<u>NAME</u>	PERCENT SUPPORTED	National Academy Member
A. P. Young	0.22	
FTE Equivalent:	0.22	
Total Number:	1	

## Names of Under Graduate students supported

<u>NAME</u>	PERCENT_SUPPORTED	
FTE Equivalent:		
Total Number:		

## **Student Metrics**

This section only applies to graduating undergraduates supported by this agreement in this reporting period

The number of undergraduates funded by this agreement who graduated during this period:	0.00
The number of undergraduates funded by this agreement who graduated during this period with a degree in	
science, mathematics, engineering, or technology fields:	. 0.00

The number of undergraduates funded by your agreement who graduated during this period and will continue to pursue a graduate or Ph.D. degree in science, mathematics, engineering, or technology fields:..... 0.00

Number of graduating undergraduates who achieved a 3.5 GPA to 4.0 (4.0 max scale):..... 0.00 Number of graduating undergraduates funded by a DoD funded Center of Excellence grant for

Education, Research and Engineering: ..... 0.00

The number of undergraduates funded by your agreement who graduated during this period and intend to work for the Department of Defense ...... 0.00

The number of undergraduates funded by your agreement who graduated during this period and will receive scholarships or fellowships for further studies in science, mathematics, engineering or technology fields: ..... 0.00

### Names of Personnel receiving masters degrees

NAME	
Total Number:	
	Names of personnel receiving PHDs
<u>NAME</u>	
Total Number:	
	Names of other research staff
<u>NAME</u>	PERCENT_SUPPORTED
FTE Equivalent: Total Number:	

**Sub Contractors (DD882)** 

**Inventions (DD882)** 

**Scientific Progress** 

**Technology Transfer** 

# Agreement Number: W911NF0910391 "Determining the Complexity of the Quantum Adiabatic Algorithm using Quantum Monte Carlo simulations" Principal Investigator: A. Peter Young

## 1 Introduction

This project is to study whether an eventual quantum computer could solve optimization problems more efficiently than a classical computer using the "quantum adiabatic algorithm" (QAA) described below.

The grant has led to six publications [1, 2, 3, 4, 5, 6], one of which [2] is a conference presentation. In addition, the PI presented results at nine international conferences which did not have proceedings. Both the PI and post-doc gave additional seminars on the research.

The initial work was done in collaboration with Smelyanskiy and Sergey Knysh, and a visiting student, Marco Guidetti participated in one of the papers, Ref. [3]. The work for the last paper [6] started during fall 2010 when the PI was on sabbatical at MIT in the group of Eddie Farhi (whose group invented the QAA in 2001). This paper [6] involved a collaboration of Farhi, his student David Gosset, Peter Shor (also at MIT), Francesco Zamponi from Paris, and Anders Sandvik from Boston University, in addition to the PI and post-doc.

Since the detailed results are in the papers, it does not seem to be necessary to repeat them all here. Rather I will give a general introduction followed by comments on the work in each of the papers. Naturally the papers also include a much larger bibliography.

For some *specialized* problems, such as factoring large integers, which is relevant for cryptography, a quantum computer is *known* to be much more efficient than a classical computer [7, 8]. The present project investigated whether, in addition, a quantum computer could solve a *wide range* of hard "optimization" problems more efficiently than a classical computer using the quantum adiabatic algorithm (QAA) [9]. The main technique used was Quantum Monte Carlo (QMC) simulations, which both the PI and the post-doc are very familiar with. Exact diagonalization was also used in some of the work. In addition, a comparison was made between the results from the QAA and a *classical* heuristic algorithm called WalkSAT.

In the QAA, a quantum Hamiltonian is realized by the connections in a quantum computer. A time-dependent control parameter, s(t),  $(0 \le s(t) \le 1)$ , is smoothly varied, taking the system from a "driver Hamiltonian"  $\mathcal{H}_{D}$ , which is easy to solve, at t = 0 (s(0) = 0), to the non-trivial "problem Hamiltonian" of interest  $\mathcal{H}_{P}$  at time  $\mathcal{T}$   $(s(\mathcal{T}) = 1)$ . The total Hamiltonian is therefore

$$\mathcal{H}(t) = [1 - s(t)]\mathcal{H}_{D} + s(t)\mathcal{H}_{P}. \tag{1}$$

For  $\mathcal{H}_{P}$  we consider binary optimization problems expressed in terms of classical bits  $z_{i} = (0,1)$ , which we represent by Ising spins taking values  $\sigma_{i}^{z} = 1 - 2z_{i} = (1,-1)$ . The simplest driver Hamiltonian is then

$$\mathcal{H}_{D} = -h \sum_{i=1}^{N} \sigma_{i}^{x}, \tag{2}$$

where  $\sigma_i^x$  is the x-component Pauli matrix and h is a "transverse field".

The quantum computer is started in the (trivial) ground state of the driver Hamiltonian and, if the process is slow enough to be adiabatic, it ends up in the ground state of the problem Hamiltonian, which solves the problem. The time taken to do this,  $\mathcal{T}$ , is called the complexity. Since the process is slow and smooth, it is easier to screen out noise, which would destroy quantum coherence, in the QAA than in other quantum algorithms, which proceed by a series of discrete unitary transformations. Hence there is a substantial experimental effort to build a quantum computer to implement the QAA. In particular, D-Wave has built a quantum annealer with over 100 qubits to implement the QAA. One of these will be installed at NASA-Ames, with financial support from Google. The PI on supported by the grant, Itay Hen, now at NASA-Ames, will provide theoretical support for this project.

An important class of optimization problems is that known as "NP-hard" [10], since their complexity is found to grow exponentially with problem size N, and, furthermore, any problem in a wide set, call "NP", can be mapped into an NP-hard problem in polynomial time.

A bottleneck will be where the gap between the ground state and the first excited state becomes very small, since the time needed to maintain adiabatic evolution is proportional to the square of the inverse of the minimum gap. The QMC simulations compute the gap as a function of the control parameter s for each instance, locate its minimum value, and see how the minimum gap, averaged in some suitable way, varies with system size.

The funding was mainly used to support a full-time post-doc, Itay Hen. Some funding was also used for summer salary for the PI, and for travel to conferences. In the first year, some computers were bought to increase the computer power available for the project. The post-doc is continuing to work on related problems as a staff scientist at NASA Ames under the supervision of Vadim Smelyanskiy.

# 1.1 Reference [1]

This paper, in collaboration with Vadim Smelyanskiy and Sergey Knysh, studied a particular constraint satisfaction problem called exact cover using quantum Monte Carlo simulations. It was found that, as the size increases, there is a first order (i.e. discontinuous) quantum phase transition during the evolution of the QAA for instances of size greater than about 100. A first order transition leads to a minimum energy gap which decreases exponentially with problem size, and hence to an exponentially large time for the QAA to solve the problem.

## 1.2 Reference [2]

This is a paper presented at the Conference on Computational Physics 2009 (CCP2009) in Kaoshing, Taiwan. CCP2009 is part of the series of conferences organized by the computational divisions of American Physical Society and the European Physical Society on Physics Computing. This conference write up is short, because of page limitations, and is a summary of Ref. [1].

## 1.3 Reference [3]

This paper was in collaboration with a visiting student, Marco Guidetti. It studied the complexity of four constraint satisfaction problems using a classical, heuristic algorithm

called WalkSAT. The aim of this study was to get results from a classical algorithm with a view to eventually comparing them with those from the quantum adiabatic algorithm. This latter stage was done in Ref. [4].

Three of the four models are in the NP-hard complexity class, and one, called 3-XORSAT, is in the P complexity class for which there exists a polynomial time algorithm. All show exponential complexity for large sizes. Curiously, the hardest problem for WalkSAT is 3-XORSAT, the one in P. Although it is not surprising that WalkSAT does not solve the this problem in polynomial time, since the polynomial time algorithm is special, and is completely unrelated to the stochastic methods in WalkSAT, it is, nonetheless, quite striking that it is actually *harder* for WalkSAT than the problems in NP. The strange result that 3-XORSAT is very "glassy" (i.e. very hard to solve by general purpose algorithms) while being easy to solve by a special algorithm, has been discussed recently to illustrate problems in a claimed proof by Deolalikar that P is not equal to NP, one of the major unsolved problems in mathematics. See for example the discussion in Refs. [11, 12].

# 1.4 Reference [4]

This paper determined the complexity of several constraint satisfaction problems using the quantum adiabatic algorithm in its simplest implementation. This was done, by studying the size dependence of the gap to the first excited state of "typical" instances using QMC simulations. The result is that, at large sizes N, the complexity increases exponentially for all models that was studied. A comparison was made with the the complexity obtained from the analogous classical algorithm WalkSAT which was found in Ref. [3]. It was seen that the harder the problem is for the classical algorithm, the harder it is also for the quantum adiabatic algorithm.

This work was selected by the American Physical Society for presentation on their web site which spotlights exceptional research:

http://physics.aps.org/synopsis-for/10.1103/PhysRevE.84.061152.

# 1.5 Reference [5]

This paper proposed a method using a quantum annealer, an analog quantum computer based on the principles of quantum adiabatic evolution, to solve the graph isomorphism problem, in which one has to determine whether two graphs are isomorphic (i.e., can be transformed into each other simply by a relabeling of the vertices). The paper demonstrated the capabilities of the method by focusing on graphs with particularly high symmetry called strongly regular graphs, though some other graphs were studied as well. This was a preliminary study so only quite small graphs were considered which could be treated by diagonalization or a related method. In future work, it would be interesting to study larger sizes using Quantum Monte Carlo simulations. The paper also showed that the method is applicable, within certain limitations, to currently available quantum hardware such as D-Wave One.

## 1.6 Reference [6]

This paper was a collaboration of the PI and post-doc with Eddie Farhi, David Gosset, Peter Shor, Francesco Zamponi and Anders Sandvik. In it we studied the performance of the quantum adiabatic algorithm on random instances of two combinatorial optimization problems, 3-regular 3-XORSAT and 3-regular max-cut (a spin glass problem). We used quantum Monte Carlo and quantum cavity methods. Using these techniques we find that the quantum adiabatic algorithm fails to solve either of these problems efficiently. For the 3-XORSAT problem it fails at the quantum phase transition where the energy gap has a minimum which is exponentially small in the problem size N. For the spin glass problem it fails beyond the quantum phase transition, i.e. in the quantum spin glass phase, where, again, the minimum gap decreases exponentially with N.

## 2 Possible future work

The work supported by the grant showed that the simplest implementation of the QAA does not seem to be more efficient for solving some classes of optimization problems than a classical algorithm. However, there is a lot of flexibility in the method. To implement the QAA one has to decide on a "path in Hamiltonian space", starting with a simple Hamiltonian for which the system can be prepared in the ground state, and ending with the specified problem Hamiltonian. The present work just took the simplest path, in which the driver Hamiltonian,  $\mathcal{H}_D$ , is a sum of transverse fields, the same on all sites, and the total Hamiltonian is a linear combination of  $\mathcal{H}_D$  and the problem Hamiltonian. In future work it would be desirable to find a more clever path in Hamiltonian space to increase the value of the minimum gap. Some ideas for that, for a problem not studied by the PI, have been proposed in Ref. [13].

It would also be useful to develop further the possible application of quantum annealing (QAA) to solve the graph isomorphism problem beyond the preliminary study in Ref. [5]. The graph isomorphism problem is interesting because, like integer factoring, it is a problem for which there is no known classical polynomial-time classical algorithm, and yet it is probably not in the NP-hard category. Since a quantum polynomial-time algorithm has been found [7] for integer factoring, it is worth investigating whether the same can be done for another problem in this category, which is intermediate between P and NP-hard, namely graph isomorphism.

In all problems discussed so far in this report, the goal has been to study the exact ground state. For many problems of practical importance, it is not necessary to find the state of lowest energy; rather, any state with energy close to that of the ground state will do. For some problems, even finding an energy close to the ground state is hard [14]. It would be worthwhile to see if the QAA could find a close-by energy more efficiently than a classical algorithm.

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